The GFtype processor

(Version 3.1, March 1991)

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The preparation of this report was supported in part by the National Science Foundation under grants IST-8201926, MCS-8300984, and CCR-8610181, and by the System Development Foundation. 'TEX' is a trademark of the American Mathematical Society. 'METAFONT' is a trademark of Addison-Wesley Publishing Company.

1. Introduction. The GFtype utility program reads binary generic-font ("GF") files that are produced by font compilers such as METAFONT, and converts them into symbolic form. This program has three chief purposes: (1) It can be used to look at the pixels of a font, with one pixel per character in a text file; (2) it can be used to determine whether a GF file is valid or invalid, when diagnosing compiler errors; and (3) it serves as an example of a program that reads GF files correctly, for system programmers who are developing GF-related software.

The original version of this program was written by David R. Fuchs in March, 1984. Donald E. Knuth made a few modifications later that year as METAFONT was taking shape.

The *banner* string defined here should be changed whenever GFtype gets modified.

define $banner \equiv \text{This}_{\sqcup}\text{is}_{\sqcup}\text{GFtype}_{,\sqcup}\text{Version}_{\sqcup}3.1^{-} \{ \text{printed when the program starts} \}$

2. This program is written in standard Pascal, except where it is necessary to use extensions; for example, one extension is to use a default **case** as in TANGLE, WEAVE, etc. All places where nonstandard constructions are used have been listed in the index under "system dependencies."

define $othercases \equiv others: { default for cases not listed explicitly }$ **define** $<math>endcases \equiv end$ { follows the default case in an extended **case** statement } **format** $othercases \equiv else$ **format** $endcases \equiv end$

3. The binary input comes from gf_{-file} , and the symbolic output is written on Pascal's standard *output* file. The term *print* is used instead of *write* when this program writes on *output*, so that all such output could easily be redirected if desired.

```
define print(#) \equiv write(#)

define print\_ln(#) \equiv write\_ln(#)

define print\_nl \equiv write\_ln

program GF\_type(gf\_file, output);

label \langle Labels in the outer block 4 \rangle

const \langle Constants in the outer block 5 \rangle

type \langle Types in the outer block 8 \rangle

var \langle Globals in the outer block 10 \rangle

procedure initialize; { this procedure gets things started properly }

var i: integer; { loop index for initializations }

begin print\_ln(banner);

\langle Set initial values 11 \rangle

end;
```

4. If the program has to stop prematurely, it goes to the 'final_end'.

define $final_end = 9999$ { label for the end of it all }

 $\label{eq:labels} \langle \, \mbox{Labels in the outer block } 4 \, \rangle \equiv $final_end$;$

This code is used in section 3.

Four parameters can be changed at compile time to extend or reduce GFtype's capacity. Note that the 5. total number of bits in the main *image_array* will be

 $(max_row + 1) \times (max_col + 1).$

(METAFONT's full pixel range is rarely implemented, because it would require 8 megabytes of memory.)

 $\langle \text{Constants in the outer block 5} \rangle \equiv$ $terminal_line_length = 150;$ { maximum number of characters input in a single line of input from the terminal } $line_length = 79; \{xxx \text{ strings will not produce lines longer than this}\}$ $max_row = 79; \{ vertical extent of pixel image array \}$ $max_col = 79; \{ \text{horizontal extent of pixel image array} \}$

This code is used in section 3.

Here are some macros for common programming idioms. 6.

define $incr(\#) \equiv \# \leftarrow \# + 1$ { increase a variable by unity } **define** $decr(\#) \equiv \# \leftarrow \# - 1$ {decrease a variable by unity } **define** $negate(\#) \equiv \# \leftarrow -\#$ { change the sign of a variable }

7. If the GF file is badly malformed, the whole process must be aborted; GFtype will give up, after issuing an error message about the symptoms that were noticed.

Such errors might be discovered inside of subroutines inside of subroutines, so a procedure called *jump_out* has been introduced. This procedure, which simply transfers control to the label *final_end* at the end of the program, contains the only non-local goto statement in GFtype.

```
define abort(\#) \equiv
               begin print(\_\_, \#); jump_out;
               end
   define bad_{-}qf(\#) \equiv abort( \mathsf{`Bad}_{||}\mathsf{GF}_{||}\mathsf{file}_{||}^{*}, \#, \mathsf{'}!^{*})
procedure jump_out;
  begin goto final_end;
   end;
```

8. The character set. Like all programs written with the WEB system, GFtype can be used with any character set. But it uses ASCII code internally, because the programming for portable input-output is easier when a fixed internal code is used.

The next few sections of GFtype have therefore been copied from the analogous ones in the WEB system routines. They have been considerably simplified, since GFtype need not deal with the controversial ASCII codes less than '40 or greater than '176. If such codes appear in the GF file, they will be printed as question marks.

 $\langle \text{Types in the outer block 8} \rangle \equiv ASCII_code = "_" ... "~"; {a subrange of the integers} See also sections 9, 20, and 36. This code is used in section 3.$

9. The original Pascal compiler was designed in the late 60s, when six-bit character sets were common, so it did not make provision for lower case letters. Nowadays, of course, we need to deal with both upper and lower case alphabets in a convenient way, especially in a program like GFtype. So we shall assume that the Pascal system being used for GFtype has a character set containing at least the standard visible characters of ASCII code ("!" through "~").

Some Pascal compilers use the original name *char* for the data type associated with the characters in text files, while other Pascals consider *char* to be a 64-element subrange of a larger data type that has some other name. In order to accommodate this difference, we shall use the name *text_char* to stand for the data type of the characters in the output file. We shall also assume that *text_char* consists of the elements $chr(first_text_char)$ through $chr(last_text_char)$, inclusive. The following definitions should be adjusted if necessary.

define $text_char \equiv char$ { the data type of characters in text files } **define** $first_text_char = 0$ { ordinal number of the smallest element of $text_char$ } **define** $last_text_char = 127$ { ordinal number of the largest element of $text_char$ }

 $\langle \text{Types in the outer block 8} \rangle +\equiv$

 $text_file =$ **packed file of** $text_char;$

10. The GFtype processor converts between ASCII code and the user's external character set by means of arrays *xord* and *xchr* that are analogous to Pascal's *ord* and *chr* functions.

 $\langle \text{Globals in the outer block 10} \rangle \equiv xord: array [text_char] of ASCII_code; { specifies conversion of input characters } xchr: array [0...255] of text_char; { specifies conversion of output characters } See also sections 21, 23, 25, 27, 35, 37, 39, 41, 46, 54, 62, and 67. This code is used in section 3.$

11. Under our assumption that the visible characters of standard ASCII are all present, the following assignment statements initialize the *xchr* array properly, without needing any system-dependent changes. $\langle \text{Set initial values } 11 \rangle \equiv$

for $i \leftarrow 0$ to '37 do $xchr[i] \leftarrow ??';$ $xchr['40] \leftarrow [']; xchr['41] \leftarrow [!]; xchr['42] \leftarrow ["]; xchr['43] \leftarrow [#]; xchr['44] \leftarrow [$;$ $xchr[45] \leftarrow \%; xchr[46] \leftarrow \&; xchr[47] \leftarrow \cdots;$ $xchr[50] \leftarrow (; xchr[51] \leftarrow); xchr[52] \leftarrow *; xchr[53] \leftarrow +; xchr[54] \leftarrow ,;$ $xchr['55] \leftarrow -; xchr['56] \leftarrow :; xchr['57] \leftarrow '/;$ $xchr[60] \leftarrow 0^{\circ}; xchr[61] \leftarrow 1^{\circ}; xchr[62] \leftarrow 2^{\circ}; xchr[63] \leftarrow 3^{\circ}; xchr[64] \leftarrow 4^{\circ};$ $xchr[65] \leftarrow 5; xchr[66] \leftarrow 6; xchr[67] \leftarrow 7;$ $xchr['70] \leftarrow \mathbf{^{*7}0} \leftarrow \mathbf{^{*7}1} \leftarrow \mathbf{^{$ $xchr[775] \leftarrow =; xchr[776] \leftarrow >; xchr[777] \leftarrow ??;$ $xchr[100] \leftarrow \text{`$@`; } xchr[101] \leftarrow \text{`$A`; } xchr[102] \leftarrow \text{`$B`; } xchr[103] \leftarrow \text{`$C`; } xchr[104] \leftarrow \text{`$D`; }$ $xchr['105] \leftarrow `E`; xchr['106] \leftarrow `F`; xchr['107] \leftarrow `G`;$ $xchr['110] \leftarrow `H`; xchr['111] \leftarrow `I`; xchr['112] \leftarrow `J`; xchr['113] \leftarrow `K`; xchr['114] \leftarrow `L`;$ $xchr['115] \leftarrow \mathsf{M}; xchr['116] \leftarrow \mathsf{N}; xchr['117] \leftarrow \mathsf{O};$ $xchr['120] \leftarrow `P`; \ xchr['121] \leftarrow `Q`; \ xchr['122] \leftarrow `R`; \ xchr['123] \leftarrow `S`; \ xchr['124] \leftarrow `T`;$ $xchr['125] \leftarrow `U`; xchr['126] \leftarrow `V`; xchr['127] \leftarrow `W`;$ $xchr['130] \leftarrow `X`; xchr['131] \leftarrow `Y`; xchr['132] \leftarrow `Z`; xchr['133] \leftarrow `['; xchr['134] \leftarrow `\`;$ $xchr['135] \leftarrow `]`; xchr['136] \leftarrow ```; xchr['137] \leftarrow `_`;$ $xchr['140] \leftarrow ```; xchr['141] \leftarrow `a`; xchr['142] \leftarrow `b`; xchr['143] \leftarrow `c`; xchr['144] \leftarrow `d`;$ $xchr['145] \leftarrow \text{`e'}; xchr['146] \leftarrow \text{`f'}; xchr['147] \leftarrow \text{`g'};$ $xchr[[150] \leftarrow `h`; xchr[[151] \leftarrow `i`; xchr[[152] \leftarrow `j`; xchr[[153] \leftarrow `k`; xchr[[154] \leftarrow `l`;$ $xchr[155] \leftarrow \text{`m'}; xchr[156] \leftarrow \text{`n'}; xchr[157] \leftarrow \text{`o'};$ $xchr['160] \leftarrow \mathbf{\hat{p}}; xchr['161] \leftarrow \mathbf{\hat{q}}; xchr['162] \leftarrow \mathbf{\hat{r}}; xchr['163] \leftarrow \mathbf{\hat{s}}; xchr['164] \leftarrow \mathbf{\hat{t}};$ $xchr['165] \leftarrow `u`; xchr['166] \leftarrow `v`; xchr['167] \leftarrow `w`;$ $xchr['170] \leftarrow \mathbf{x}; xchr['171] \leftarrow \mathbf{y}; xchr['172] \leftarrow \mathbf{z}; xchr['173] \leftarrow \mathbf{f}; xchr['174] \leftarrow \mathbf{f};$ $xchr['175] \leftarrow ``` ; xchr['176] \leftarrow ```$ for $i \leftarrow 177$ to 255 do $xchr[i] \leftarrow ??$;

See also sections 12, 26, 47, and 63.

This code is used in section **3**.

12. The following system-independent code makes the *xord* array contain a suitable inverse to the information in *xchr*.

 $\langle \text{Set initial values } 11 \rangle + \equiv$

for $i \leftarrow first_text_char$ to $last_text_char$ do $xord[chr(i)] \leftarrow '40$; for $i \leftarrow "_$ " to "~" do $xord[xchr[i]] \leftarrow i$; 13. Generic font file format. The most important output produced by a typical run of METAFONT is the "generic font" (GF) file that specifies the bit patterns of the characters that have been drawn. The term generic indicates that this file format doesn't match the conventions of any name-brand manufacturer; but it is easy to convert GF files to the special format required by almost all digital phototypesetting equipment. There's a strong analogy between the DVI files written by TEX and the GF files written by METAFONT; and, in fact, the file formats have a lot in common. It is therefore not surprising that GFtype is identical in many respects to the DVItype program.

A GF file is a stream of 8-bit bytes that may be regarded as a series of commands in a machine-like language. The first byte of each command is the operation code, and this code is followed by zero or more bytes that provide parameters to the command. The parameters themselves may consist of several consecutive bytes; for example, the 'boc' (beginning of character) command has six parameters, each of which is four bytes long. Parameters are usually regarded as nonnegative integers; but four-byte-long parameters can be either positive or negative, hence they range in value from -2^{31} to $2^{31} - 1$. As in TFM files, numbers that occupy more than one byte position appear in BigEndian order, and negative numbers appear in two's complement notation.

A GF file consists of a "preamble," followed by a sequence of one or more "characters," followed by a "postamble." The preamble is simply a *pre* command, with its parameters that introduce the file; this must come first. Each "character" consists of a *boc* command, followed by any number of other commands that specify "black" pixels, followed by an *eoc* command. The characters appear in the order that METAFONT generated them. If we ignore no-op commands (which are allowed between any two commands in the file), each *eoc* command is immediately followed by a *boc* command, or by a *post* command; in the latter case, there are no more characters in the file, and the remaining bytes form the postamble. Further details about the postamble will be explained later.

Some parameters in GF commands are "pointers." These are four-byte quantities that give the location number of some other byte in the file; the first file byte is number 0, then comes number 1, and so on.

14. The GF format is intended to be both compact and easily interpreted by a machine. Compactness is achieved by making most of the information relative instead of absolute. When a GF-reading program reads the commands for a character, it keeps track of two quantities: (a) the current column number, m; and (b) the current row number, n. These are 32-bit signed integers, although most actual font formats produced from GF files will need to curtail this vast range because of practical limitations. (METAFONT output will never allow |m| or |n| to get extremely large, but the GF format tries to be more general.)

How do GF's row and column numbers correspond to the conventions of T_EX and METAFONT? Well, the "reference point" of a character, in T_EX's view, is considered to be at the lower left corner of the pixel in row 0 and column 0. This point is the intersection of the baseline with the left edge of the type; it corresponds to location (0,0) in METAFONT programs. Thus the pixel in GF row 0 and column 0 is METAFONT's unit square, comprising the region of the plane whose coordinates both lie between 0 and 1. The pixel in GF row n and column m consists of the points whose METAFONT coordinates (x, y) satisfy $m \le x \le m + 1$ and $n \le y \le n + 1$. Negative values of m and x correspond to columns of pixels left of the reference point; negative values of n and y correspond to rows of pixels below the baseline.

Besides m and n, there's also a third aspect of the current state, namely the *paint_switch*, which is always either *black* or *white*. Each *paint* command advances m by a specified amount d, and blackens the intervening pixels if *paint_switch* = *black*; then the *paint_switch* changes to the opposite state. GF's commands are designed so that m will never decrease within a row, and n will never increase within a character; hence there is no way to whiten a pixel that has been blackened. 15. Here is a list of all the commands that may appear in a GF file. Each command is specified by its symbolic name (e.g., *boc*), its opcode byte (e.g., 67), and its parameters (if any). The parameters are followed by a bracketed number telling how many bytes they occupy; for example, d[2] means that parameter d is two bytes long.

- $paint_0 0$. This is a paint command with d = 0; it does nothing but change the paint_switch from black to white or vice versa.
- paint_1 through paint_63 (opcodes 1 to 63). These are paint commands with d = 1 to 63, defined as follows: If paint_switch = black, blacken d pixels of the current row n, in columns m through m + d - 1 inclusive. Then, in any case, complement the paint_switch and advance m by d.
- paint1 64 d[1]. This is a paint command with a specified value of d; METAFONT uses it to paint when $64 \le d < 256$.
- paint2 65 d[2]. Same as paint1, but d can be as high as 65535.
- paint3 66 d[3]. Same as paint1, but d can be as high as $2^{24} 1$. METAFONT never needs this command, and it is hard to imagine anybody making practical use of it; surely a more compact encoding will be desirable when characters can be this large. But the command is there, anyway, just in case.
- boc 67 c[4] p[4] $min_m[4]$ $max_m[4]$ $min_n[4]$ $max_n[4]$. Beginning of a character: Here c is the character code, and p points to the previous character beginning (if any) for characters having this code number modulo 256. (The pointer p is -1 if there was no prior character with an equivalent code.) The values of registers m and n defined by the instructions that follow for this character must satisfy $min_m \leq m \leq max_m$ and $min_n \leq n \leq max_n$. (The values of max_m and min_n need not be the tightest bounds possible.) When a GF-reading program sees a boc, it can use min_m , max_m , min_n , and max_n to initialize the bounds of an array. Then it sets $m \leftarrow min_m$, $n \leftarrow max_n$, and $paint_switch \leftarrow white$.
- boc1 68 c[1] $del_m[1] max_m[1] del_n[1] max_n[1]$. Same as boc, but p is assumed to be -1; also $del_m = max_m min_m$ and $del_n = max_n min_n$ are given instead of min_m and min_n . The one-byte parameters must be between 0 and 255, inclusive. (This abbreviated boc saves 19 bytes per character, in common cases.)
- eoc 69. End of character: All pixels blackened so far constitute the pattern for this character. In particular, a completely blank character might have eoc immediately following boc.
- skip0 70. Decrease n by 1 and set $m \leftarrow min_m$, $paint_switch \leftarrow white$. (This finishes one row and begins another, ready to whiten the leftmost pixel in the new row.)
- skip1 71 d[1]. Decrease n by d+1, set $m \leftarrow min_m$, and set paint_switch \leftarrow white. This is a way to produce d all-white rows.
- skip2 72 d[2]. Same as skip1, but d can be as large as 65535.
- skip3 73 d[3]. Same as skip1, but d can be as large as $2^{24} 1$. METAFONT obviously never needs this command.
- new_row_0 74. Decrease n by 1 and set $m \leftarrow min_m$, $paint_switch \leftarrow black$. (This finishes one row and begins another, ready to blacken the leftmost pixel in the new row.)
- *new_row_1* through *new_row_164* (opcodes 75 to 238). Same as *new_row_0*, but with $m \leftarrow min_m + 1$ through $min_m + 164$, respectively.
- xxx1 239 k[1] x[k]. This command is undefined in general; it functions as a (k + 2)-byte no_op unless special GF-reading programs are being used. METAFONT generates xxx commands when encountering a **special** string; this occurs in the GF file only between characters, after the preamble, and before the postamble. However, xxx commands might appear within characters, in GF files generated by other processors. It is recommended that x be a string having the form of a keyword followed by possible parameters relevant to that keyword.
- *xxx2* 240 k[2] x[k]. Like *xxx1*, but $0 \le k < 65536$.
- xxx3 241 k[3] x[k]. Like xxx1, but $0 \le k < 2^{24}$. METAFONT uses this when sending a special string whose length exceeds 255.

xxx4 242 k[4] x[k]. Like xxx1, but k can be ridiculously large; k mustn't be negative.

- yyy 243 y[4]. This command is undefined in general; it functions as a 5-byte no_op unless special GF-reading programs are being used. METAFONT puts *scaled* numbers into yyy's, as a result of **numspecial** commands; the intent is to provide numeric parameters to xxx commands that immediately precede.
- *no_op* 244. No operation, do nothing. Any number of *no_op*'s may occur between **GF** commands, but a *no_op* cannot be inserted between a command and its parameters or between two parameters.
- char_loc 245 c[1] dx[4] dy[4] w[4] p[4]. This command will appear only in the postamble, which will be explained shortly.
- char_loc0 246 c[1] dm[1] w[4] p[4]. Same as char_loc, except that dy is assumed to be zero, and the value of dx is taken to be 65536 * dm, where $0 \le dm < 256$.
- pre 247 i[1] k[1] x[k]. Beginning of the preamble; this must come at the very beginning of the file. Parameter i is an identifying number for GF format, currently 131. The other information is merely commentary; it is not given special interpretation like xxx commands are. (Note that xxx commands may immediately follow the preamble, before the first *boc*.)
- post 248. Beginning of the postamble, see below.

post_post 249. Ending of the postamble, see below.

Commands 250–255 are undefined at the present time.

define $gf_id_byte = 131$ {identifies the kind of GF files described here}

16. Here are the opcodes that GFtype actually refers to.

define $paint_0 = 0$ { beginning of the paint commands } **define** paint1 = 64 {move right a given number of columns, then black \leftrightarrow white } **define** boc = 67 { beginning of a character } **define** boc1 = 68 { abbreviated boc } **define** eoc = 69 { end of a character } **define** $skip\theta = 70 \{ skip no blank rows \}$ **define** skip1 = 71 { skip over blank rows } **define** $new_row_0 = 74$ { move down one row and then right } define xxx1 = 239 { for special strings } define yyy = 243 { for numspecial numbers } define $no_op = 244 \{ no operation \}$ define $char_{loc} = 245$ { character locators in the postamble } define $pre = 247 \{ \text{preamble} \}$ **define** post = 248 { postamble beginning } **define** $post_post = 249$ { postamble ending } **define** $undefined_commands \equiv 250, 251, 252, 253, 254, 255$

17. The last character in a GF file is followed by '*post*'; this command introduces the postamble, which summarizes important facts that METAFONT has accumulated. The postamble has the form

post p[4] ds[4] cs[4] hppp[4] vppp[4] min_m[4] max_m[4] min_n[4] max_n[4] (character locators) post_post q[4] i[1] 223's[\geq 4]

Here p is a pointer to the byte following the final *eoc* in the file (or to the byte following the preamble, if there are no characters); it can be used to locate the beginning of xxx commands that might have preceded the postamble. The ds and cs parameters give the design size and check sum, respectively, which are exactly the values put into the header of any TFM file that shares information with this GF file. Parameters hppp and vppp are the ratios of pixels per point, horizontally and vertically, expressed as scaled integers (i.e., multiplied by 2^{16}); they can be used to correlate the font with specific device resolutions, magnifications, and "at sizes." Then come min_m, max_m, min_n, and max_n, which bound the values that registers m and n assume in all characters in this GF file. (These bounds need not be the best possible; max_m and min_n may, on the other hand, be tighter than the similar bounds in boc commands. For example, some character may have min_n = -100 in its boc, but it might turn out that n never gets lower than -50 in any character; then min_n can have any value ≤ -50 . If there are no characters in the file, it's possible to have min_m > max_m and/or min_n > max_n.)

18. Character locators are introduced by $char_loc$ commands, which specify a character residue c, character escapements (dx, dy), a character width w, and a pointer p to the beginning of that character. (If two or more characters have the same code c modulo 256, only the last will be indicated; the others can be located by following backpointers. Characters whose codes differ by a multiple of 256 are assumed to share the same font metric information, hence the TFM file contains only residues of character codes modulo 256. This convention is intended for oriental languages, when there are many character shapes but few distinct widths.)

The character escapements (dx, dy) are the values of METAFONT's **chardx** and **chardy** parameters; they are in units of *scaled* pixels; i.e., dx is in horizontal pixel units times 2^{16} , and dy is in vertical pixel units times 2^{16} . This is the intended amount of displacement after typesetting the character; for DVI files, dy should be zero, but other document file formats allow nonzero vertical escapement.

The character width w duplicates the information in the TFM file; it is 2^{24} times the ratio of the true width to the font's design size.

The backpointer p points to the character's *boc*, or to the first of a sequence of consecutive xxx or yyy or no_op commands that immediately precede the *boc*, if such commands exist; such "special" commands essentially belong to the characters, while the special commands after the final character belong to the postamble (i.e., to the font as a whole). This convention about p applies also to the backpointers in *boc* commands, even though it wasn't explained in the description of *boc*.

Pointer p might be -1 if the character exists in the TFM file but not in the GF file. This unusual situation can arise in METAFONT output if the user had proofing < 0 when the character was being shipped out, but then made proofing ≥ 0 in order to get a GF file.

19. The last part of the postamble, following the *post_post* byte that signifies the end of the character locators, contains q, a pointer to the *post* command that started the postamble. An identification byte, i, comes next; this currently equals 131, as in the preamble.

The *i* byte is followed by four or more bytes that are all equal to the decimal number 223 (i.e., "DF in hexadecimal). METAFONT puts out four to seven of these trailing bytes, until the total length of the file is a multiple of four bytes, since this works out best on machines that pack four bytes per word; but any number of 223's is allowed, as long as there are at least four of them. In effect, 223 is a sort of signature that is added at the very end.

This curious way to finish off a GF file makes it feasible for GF-reading programs to find the postamble first, on most computers, even though METAFONT wants to write the postamble last. Most operating systems permit random access to individual words or bytes of a file, so the GF reader can start at the end and skip backwards over the 223's until finding the identification byte. Then it can back up four bytes, read q, and move to byte q of the file. This byte should, of course, contain the value 248 (*post*); now the postamble can be read, so the GF reader can discover all the information needed for individual characters.

Unfortunately, however, standard Pascal does not include the ability to access a random position in a file, or even to determine the length of a file. Almost all systems nowadays provide the necessary capabilities, so GF format has been designed to work most efficiently with modern operating systems. But if GF files have to be processed under the restrictions of standard Pascal, one can simply read them from front to back. This will be adequate for most applications. However, the postamble-first approach would facilitate a program that merges two GF files, replacing data from one that is overridden by corresponding data in the other.

20. Input from binary files. We have seen that a GF file is a sequence of 8-bit bytes. The bytes appear physically in what is called a 'packed file of 0..255' in Pascal lingo.

Packing is system dependent, and many Pascal systems fail to implement such files in a sensible way (at least, from the viewpoint of producing good production software). For example, some systems treat all byte-oriented files as text, looking for end-of-line marks and such things. Therefore some system-dependent code is often needed to deal with binary files, even though most of the program in this section of GFtype is written in standard Pascal.

We shall stick to simple Pascal in this program, for reasons of clarity, even if such simplicity is sometimes unrealistic.

 $\langle \text{Types in the outer block 8} \rangle +\equiv eight_bits = 0..255;$ { unsigned one-byte quantity } byte_file = packed file of eight_bits; { files that contain binary data }

21. The program deals with one binary file variable: gf_{-file} is the main input file that we are translating into symbolic form.

 $\langle \text{Globals in the outer block 10} \rangle +\equiv gf_file: byte_file; { the stuff we are GFtyping }$

22. To prepare this file for input, we reset it.

procedure *open_gf_file*; { prepares to read packed bytes in *gf_file* } **begin** *reset*(*gf_file*); *cur_loc* \leftarrow 0; **end**;

23. If you looked carefully at the preceding code, you probably asked, "What is *cur_loc*?" Good question. It's a global variable that holds the number of the byte about to be read next from *gf_file*.

 $\langle \text{Globals in the outer block } 10 \rangle + \equiv$

cur_loc: integer; { where we are about to look, in gf_file }

24. We shall use a set of simple functions to read the next byte or bytes from $gf_{-}file$. There are four possibilities, each of which is treated as a separate function in order to minimize the overhead for subroutine calls.

function *get_byte*: *integer*; { returns the next byte, unsigned } **var** b: eight_bits; **begin if** $eof(gf_file)$ **then** $get_byte \leftarrow 0$ else begin read (gf_file, b); incr(cur_loc); get_byte \leftarrow b; end; end; **function** get_two_bytes: integer; { returns the next two bytes, unsigned } **var** a, b: eight_bits; **begin** read (gf_file, a) ; read (gf_file, b) ; cur_loc \leftarrow cur_loc + 2; get_two_bytes $\leftarrow a * 256 + b$; end; function get_three_bytes: integer; { returns the next three bytes, unsigned } **var** a, b, c: eight_bits; **begin** read (gf_{file}, a); read (gf_{file}, b); read (gf_{file}, c); cur_loc \leftarrow cur_loc + 3; get_three_bytes $\leftarrow (a * 256 + b) * 256 + c;$ end; function *signed_quad*: *integer*; { returns the next four bytes, signed } **var** a, b, c, d: *eight_bits*; **begin** read (qf_{file}, a) ; read (qf_{file}, b) ; read (qf_{file}, c) ; read (qf_{file}, d) ; cur_loc \leftarrow cur_loc + 4; if a < 128 then signed_quad $\leftarrow ((a * 256 + b) * 256 + c) * 256 + d$ else $signed_quad \leftarrow (((a - 256) * 256 + b) * 256 + c) * 256 + d;$ end;

25. Optional modes of output. GFtype will print different quantities of information based on some options that the user must specify: We set *wants_mnemonics* if the user wants to see a mnemonic dump of the GF file; and we set *wants_pixels* if the user wants to see a pixel image of each character.

When GFtype begins, it engages the user in a brief dialog so that the options will be specified. This part of GFtype requires nonstandard Pascal constructions to handle the online interaction; so it may be preferable in some cases to omit the dialog and simply to produce the maximum possible output (wants_mnemonics = wants_pixels = true). On other hand, the necessary system-dependent routines are not complicated, so they can be introduced without terrible trauma.

 $\langle \text{Globals in the outer block } 10 \rangle + \equiv$ wants_mnemonics: boolean; { controls mnemonic output } wants_pixels: boolean; { controls pixel output }

```
26. \langle Set initial values 11 \rangle +\equiv
wants_mnemonics \leftarrow true; wants_pixels \leftarrow true;
```

27. The *input_ln* routine waits for the user to type a line at his or her terminal; then it puts ASCII-code equivalents for the characters on that line into the *buffer* array. The *term_in* file is used for terminal input, and *term_out* for terminal output.

 $\langle \text{Globals in the outer block } 10 \rangle +\equiv buffer: array [0..terminal_line_length] of ASCII_code; term_in: text_file; { the terminal, considered as an input file } term_out: text_file; { the terminal, considered as an output file }$

28. Since the terminal is being used for both input and output, some systems need a special routine to make sure that the user can see a prompt message before waiting for input based on that message. (Otherwise the message may just be sitting in a hidden buffer somewhere, and the user will have no idea what the program is waiting for.) We shall invoke a system-dependent subroutine *update_terminal* in order to avoid this problem.

define $update_terminal \equiv break(term_out)$ { empty the terminal output buffer }

29. During the dialog, extensions of GFtype might treat the first blank space in a line as the end of that line. Therefore *input_ln* makes sure that there is always at least one blank space in *buffer*.

(This routine is more complex than the present implementation needs, but it has been copied from DVItype so that system-dependent changes that worked before will work again.)

procedure *input_ln*; { inputs a line from the terminal }

```
var k: 0 .. terminal_line_length;
begin update_terminal; reset(term_in);
if eoln(term_in) then read_ln(term_in);
k \leftarrow 0;
while (k < terminal_line_length) \land \neg eoln(term_in) do
begin buffer[k] \leftarrow word[term_in^{\uparrow}]; incr(k); get(term_in);
end;
buffer[k] \leftarrow "_";
end;
```

30. This is humdrum.

```
function lower\_casify(c: ASCII\_code): ASCII\_code;
begin if (c \ge "A") \land (c \le "Z") then lower\_casify \leftarrow c + "a" - "A"
else lower\_casify \leftarrow c;
end;
```

31. The selected options are put into global variables by the *dialog* procedure, which is called just as GFtype begins.

procedure dialog;

32. (Determine whether the user *wants_mnemonics* 32) \equiv

1: write(term_out, `Mnemonic_output?_(default=no,_?_for_help):_`); input_ln; buffer[0] ← lower_casify(buffer[0]); if buffer[0] ≠ "?" then wants_mnemonics ← (buffer[0] = "y") ∨ (buffer[0] = "1") ∨ (buffer[0] = "t") else begin write(term_out, `Type_Y_for_complete_listing, `); write_ln(term_out, `_N_for_errors/images_only.`); goto 1; end

This code is used in section 31.

33. (Determine whether the user *wants_pixels* 33) \equiv

2: write(term_out, `Pixel_output?_(default=yes, _?_for_help):_'; input_ln; buffer[0] ← lower_casify(buffer[0]); if buffer[0] ≠ "?" then wants_pixels ← (buffer[0] = "y") ∨ (buffer[0] = "1") ∨ (buffer[0] = "t") ∨ (buffer[0] = "_"") else begin write(term_out, `Type_Y_ito_list_characters_pictorially`); write_ln(term_out, `_with_*``s, _N_ito_omit_this_option.`); goto 2; end

This code is used in section 31.

34. After the dialog is over, we print the options so that the user can see what GFtype thought was specified.

```
(Print all the selected options 34) =
print(`Options_selected:_Mnemonic_output_=_`);
if wants_mnemonics then print(`true`) else print(`false`);
print(`;_pixel_output_=_`);
if wants_pixels then print(`true`) else print(`false`);
print_ln(`.`)
```

This code is used in section 31.

35. The image array. The definition of GF files refers to two registers, m and n, which hold integer column and row numbers. We actually keep the values $m' = m - min_m$ and $n' = max_n - n$ instead, so that our internal image array always has $m, n \ge 0$. We also need to remember *paint_switch*, whose value is either *black* or *white*.

 $\langle \text{Globals in the outer block 10} \rangle +\equiv m, n: integer; { current state values, modified by min_m and max_n } paint_switch: pixel;$

36. We'll need a big array of pixels to hold the character image. Each pixel should be represented as a single bit in order to save space. Some systems may prefer the following definitions, while others may do better using the *boolean* type and boolean constants.

define white = 0 { could also be false } define black = 1 { could also be true } $\langle \text{Types in the outer block 8} \rangle +\equiv$ $pixel = white \dots black;$ { could also be boolean }

37. In order to allow different systems to change the *image* array easily from row-major order to columnmajor order (or vice versa), or to transpose it top and bottom or left and right, we declare and access it as follows.

define $image \equiv image_array[m, n]$ \langle Globals in the outer block 10 $\rangle +\equiv$ $image_array$: packed array [0 ... max_col, 0 ... max_row] of pixel;

38. A boc command has parameters min_m , max_m , min_n , and max_n that define a rectangular subarray in which the pixels of the current character must lie. The program here computes limits on GFtype's modified m and n variables, and clears the resulting subarray to all white.

(There may be a faster way to clear a subarray on particular systems, using nonstandard extensions of Pascal.)

```
\langle \text{Clear the image } 38 \rangle \equiv

begin max\_subcol \leftarrow max\_m\_stated - min\_m\_stated - 1;

if max\_subcol > max\_col then max\_subcol \leftarrow max\_col;

max\_subrow \leftarrow max\_n\_stated - min\_n\_stated;

if max\_subrow \leftarrow max\_n\_stated - min\_n\_stated;

if max\_subrow > max\_row then max\_subrow \leftarrow max\_row;

n \leftarrow 0;

while n \le max\_subrow do

begin m \leftarrow 0;

while m \le max\_subcol do

begin image \leftarrow white; incr(m);

end;

incr(n);

end;

thend

This code is used in section 71.
```

39. \langle Globals in the outer block $10 \rangle +\equiv max_subrow, max_subcol: integer; { size of current subarray of interest }$

40. As we paint the pixels of a character, we will record its actual boundaries in variables $max_m_observed$ and $max_n_observed$. Then the following routine will be called on to output the image, using blanks for white and asterisks for black. Blanks are emitted only when they are followed by nonblanks, in order to conserve space in the output. Further compaction could be achieved on many systems by using tab marks.

An integer variable b will be declared for use in counting blanks.

 $\langle \text{Print the image 40} \rangle \equiv$

begin \langle Compare the subarray boundaries with the observed boundaries $42\rangle$;

if $max_subcol \ge 0$ then { there was at least one paint command }

 \langle Print asterisk patterns for rows 0 to max_subrow 43 \rangle

else $print_ln((The_character_is_entirely_blank.));$

end

This code is used in section 69.

41. \langle Globals in the outer block $10 \rangle +\equiv$

 $min_m_stated, max_m_stated, min_n_stated, max_n_stated: integer;$ { bounds stated in the GF file } $max_m_observed, max_n_observed: integer;$ { bounds on (m', n') actually observed when painting } $min_m_overall, max_m_overall, min_n_overall: integer;$

{ bounds observed in the entire file so far }

42. If the given character is substantially smaller than the *boc* command predicted, we don't want to bother to output rows and columns that are all blank.

 \langle Compare the subarray boundaries with the observed boundaries $42 \rangle \equiv$

if (max_m_observed > max_col) ∨ (max_n_observed > max_row) then
 print_ln(`(The_character_is_too_large_to_be_displayed_in_full.)`);
if man or books > man m observed then man or books.

 $\label{eq:max_subcol} \mbox{if } max_subcol > max_m_observed \ \mbox{then} \ \ max_subcol \leftarrow max_m_observed;$

 $\mbox{if } max_subrow > max_n_observed \ \mbox{then} \ \ max_subrow \leftarrow max_n_observed; \\ \mbox{is } max_n_observed \ \mbox{is } max_n_observ$

This code is used in section 40.

```
43.
      \langle Print asterisk patterns for rows 0 to max_subrow 43 \rangle \equiv
  begin print_ln(`.<--This_pixel``s_lower_left_corner_is_at_(`, min_m_stated : 1, `,`,
       max_n_stated + 1: 1, `)_1 in_1 METAFONT_1 coordinates`); n \leftarrow 0;
  while n < max_{subrow} do
    begin m \leftarrow 0; b \leftarrow 0;
    while m \leq max\_subcol do
      begin if image = white then incr(b)
      else begin while b > 0 do
           begin print(\_\_]; decr(b);
           end;
         print(`*`);
         end;
       incr(m);
      end;
    print_nl; incr(n);
    end;
  print_ln(`.<--This_pixel``s_upper_left_corner_is_at_(`, min_m_stated : 1, `, `,
       max_n_stated - max_subrow : 1, `) _in_METAFONT_coordinates `);
  end
```

This code is used in section 40.

44. Translation to symbolic form. The main work of GFtype is accomplished by the *do_char* procedure, which produces the output for an entire character, assuming that the *boc* command for that character has already been processed. This procedure is essentially an interpretive routine that reads and acts on the GF commands.

45. We steal the following routine from METAFONT.

```
define unity \equiv 200000 \{2^{16}, \text{ represents } 1.00000\}

procedure print\_scaled(s:integer); \{ \text{ prints a scaled number, rounded to five digits } 

var delta: integer; \{ \text{ amount of allowable inaccuracy } \}

begin if s < 0 then

begin print(`-`); negate(s); \{ \text{ print the sign, if negative } \}

end;

print(s \operatorname{div} unity : 1); \{ \text{ print the integer part } \}

s \leftarrow 10 * (s \operatorname{mod} unity) + 5;

if s \neq 5 then

begin delta \leftarrow 10; print(`.`);

repeat if delta > unity then s \leftarrow s + 100000 - (delta \operatorname{div} 2); \{ \text{ round the final digit } \}

print(chr(ord(`0`) + (s \operatorname{div} unity))); s \leftarrow 10 * (s \operatorname{mod} unity); delta \leftarrow delta * 10;

until s \leq delta;

end;

end;
```

46. Let's keep track of how many characters are in the font, and the locations of where each one occurred in the file.

```
\langle \text{Globals in the outer block 10} \rangle +\equiv total_chars: integer; { the total number of characters seen so far } 
char_ptr: array [0...255] of integer; { correct character location pointer } 
gf_prev_ptr: integer; { char_ptr for next character } 
character_code: integer; { current character number }
```

```
47. \langle \text{Set initial values } 11 \rangle + \equiv
for i \leftarrow 0 to 255 do char_ptr[i] \leftarrow -1; { mark characters as not being in the file }
total_chars \leftarrow 0;
```

§48 GFtype

48. Before we get into the details of *do_char*, it is convenient to consider a simpler routine that computes the first parameter of each opcode.

 $\begin{array}{l} \textbf{define} \ four_cases(\texttt{\#}) \equiv \texttt{\#}, \texttt{\#} + 1, \texttt{\#} + 2, \texttt{\#} + 3\\ \textbf{define} \ eight_cases(\texttt{\#}) \equiv four_cases(\texttt{\#}), four_cases(\texttt{\#} + 4)\\ \textbf{define} \ sixteen_cases(\texttt{\#}) \equiv eight_cases(\texttt{\#}), eight_cases(\texttt{\#} + 8)\\ \textbf{define} \ thirty_two_cases(\texttt{\#}) \equiv sixteen_cases(\texttt{\#}), sixteen_cases(\texttt{\#} + 16)\\ \textbf{define} \ thirty_seven_cases(\texttt{\#}) \equiv thirty_two_cases(\texttt{\#}), four_cases(\texttt{\#} + 32), \texttt{\#} + 36\\ \textbf{define} \ sixty_four_cases(\texttt{\#}) \equiv thirty_two_cases(\texttt{\#}), thirty_two_cases(\texttt{\#} + 32)\\ \end{array}$

function *first_par(o : eight_bits): integer;*

begin case o of

```
\begin{aligned} sixty\_four\_cases(paint\_0): \ first\_par \leftarrow o - paint\_0; \\ paint1, skip1, char\_loc, char\_loc + 1, xxx1: \ first\_par \leftarrow get\_byte; \\ paint1 + 1, skip1 + 1, xxx1 + 1: \ first\_par \leftarrow get\_two\_bytes; \\ paint1 + 2, skip1 + 2, xxx1 + 2: \ first\_par \leftarrow get\_three\_bytes; \\ xxx1 + 3, yyy: \ first\_par \leftarrow signed\_quad; \\ boc, boc1, eoc, skip0, no\_op, pre, post, post\_post, undefined\_commands: \ first\_par \leftarrow 0; \\ sixty\_four\_cases(new\_row\_0), sixty\_four\_cases(new\_row\_0 + 64), thirty\_seven\_cases(new\_row\_0 + 128): \\ \ first\_par \leftarrow o - new\_row\_0; \\ end; \\ end; \end{aligned}
```

49. Strictly speaking, the *do_char* procedure is really a function with side effects, not a '**procedure**'; it returns the value *false* if **GFtype** should be aborted because of some unusual happening. The subroutine is organized as a typical interpreter, with a multiway branch on the command code.

function do_char: boolean;

```
label 9998, 9999;
var o: eight_bits; { operation code of the current command }
    p, q: integer; { parameters of the current command }
    aok: boolean; { the value to return }
    begin { we've already scanned the boc }
    aok ← true;
    while true do 〈 Translate the next command in the GF file; goto 9999 if it was eoc; goto 9998 if
        premature termination is needed 50 ⟩;
    9998: print_ln(`!`); aok ← false;
    9999: do_char ← aok;
```

end;

 \langle Translate the next command in the GF file; goto 9999 if it was *eoc*; goto 9998 if premature termination is needed 50 $\rangle \equiv$

begin *start_op*; \langle Start translation of command *o* and **goto** the appropriate label to finish the job 51 \rangle ; end

This code is used in section 49.

51. The multiway switch in *first_par*, above, was organized by the length of each command; the one in *do_char* is organized by the semantics.

 \langle Start translation of command *o* and **goto** the appropriate label to finish the job 51 $\rangle \equiv$

if $o \le paint1 + 3$ then $\langle \text{Translate a sequence of } paint \text{ commands, until reaching a non-paint } 56 \rangle$; case o of

 $four_cases(skip0)$: (Translate a skip command 60);

 $sixty_four_cases(new_row_0), sixty_four_cases(new_row_0 + 64), thirty_seven_cases(new_row_0 + 128):$ (Translate a new_row command 59);

 $\langle \text{Cases for commands } no_op, pre, post, post_post, boc, and eoc 52 \rangle$

 $four_cases(xxx1)$: (Translate an xxx command 53);

yyy: $\langle \text{Translate a yyy command 55} \rangle$;

othercases *error*(`undefined_command_`, o: 1, `!`)

endcases

This code is used in section 50.

```
52. (Cases for commands no_op, pre, post, post_post, boc, and eoc 52) \equiv
```

no_op: *show_mnemonic*(`no⊔op`);

```
pre: begin error(`preamble_command_within_a_character!`); goto 9998;
end;
```

- post, post_post: begin error(`postamble_command_within_a_character!`); goto 9998; end;
- boc, boc1: begin error(`boc_occurred_before_eoc!`); goto 9998; end:
- eoc: begin show_mnemonic(`eoc`); print_nl; goto 9999;
 end;

This code is used in section 51.

```
\langle \text{Translate an } xxx \text{ command } 53 \rangle \equiv
53.
  begin show_mnemonic(xxx_{\perp}); bad_char \leftarrow false; b \leftarrow 16;
  if p < 0 then nl_{error}(\text{`string_of_negative_length!'});
  while p > 0 do
     begin q \leftarrow get_byte;
     if (q < " \sqcup ") \lor (q > " ~ ") then bad_char \leftarrow true;
     if wants_mnemonics then
       begin print(xchr[q]);
       if b < line_length then incr(b)
       else begin print_nl; b \leftarrow 2;
          end;
       end;
     decr(p);
     end:
  if wants_mnemonics then print(````);
  if bad_char then nl_error(`non-ASCII_character_in_xxx_command!`);
  end
```

This code is used in sections 51 and 70.

54. (Globals in the outer block 10) += bad_char: boolean; { has a non-ASCII character code appeared in this xxx? }

```
55. (Translate a yyy command 55) ≡
begin show_mnemonic(`yyyu`, p: 1, `u(`);
if wants_mnemonics then
    begin print_scaled(p); print(`)`);
end;
end
```

```
This code is used in sections 51 and 70.
```

56. The bulk of a GF file generally consists of *paint* commands, so we collect them together and print them in an abbreviated format on one line.

```
( Translate a sequence of paint commands, until reaching a non-paint 56 ) ≡
begin if wants_mnemonics then print(`__paint__`);
repeat ( Paint the next p pixels 57 );
start_op;
until o > paint1 + 3;
end
This code is used in section 51.
```

57. \langle Paint the next p pixels $57 \rangle \equiv$

 $\mathbf{if} \ wants_mnemonics \ \mathbf{then}$

if $paint_switch = white then print(`(`, p:1, `)`) else print(p:1);$

 $m \leftarrow m + p;$

if $m > max_m$ -observed then max_m -observed $\leftarrow m - 1$;

if wants_pixels then (Paint pixels m - p through m - 1 in row n of the subarray 58);

 $paint_switch \leftarrow white + black - paint_switch$ { could also be $paint_switch \leftarrow \neg paint_switch$ } This code is used in section 56. 58. We use the fact that the subarray has been initialized to all *white*.

```
\langle \text{Paint pixels } m - p \text{ through } m - 1 \text{ in row } n \text{ of the subarray } 58 \rangle \equiv

if paint\_switch = black then

if n \leq max\_subrow then

begin l \leftarrow m - p; r \leftarrow m - 1;

if r > max\_subcol then r \leftarrow max\_subcol;

m \leftarrow l;

while m \leq r do

begin image \leftarrow black; incr(m);

end;

m \leftarrow l + p;

end

This code is used in section 57.
```

59. (Translate a new_row command 59) ≡ begin show_mnemonic(`newrow」`, p:1); incr(n); m ← p; paint_switch ← black; if wants_mnemonics then print(`_(n=`, max_n_stated - n:1, `)`); end

This code is used in section 51.

60. (Translate a skip command 60) ≡ begin show_mnemonic(`skip`, (o - skip1 + 1) mod 4: 1, `__`, p: 1); n ← n + p + 1; m ← 0; paint_switch ← white; if wants_mnemonics then print(`__(n=`, max_n_stated - n: 1, `)`); end

This code is used in section 51.

61. Reading the postamble. Now imagine that we are reading the GF file and positioned just after the *post* command. That, in fact, is the situation, when the following part of GFtype is called upon to read, translate, and check the rest of the postamble.

procedure read_postamble;

var k: integer; { loop index } p, q, m, u, v, w, c: integer; { general purpose registers } **begin** $post_loc \leftarrow cur_loc - 1$; $print(`Postamble_starts_at_byte_`, post_loc : 1)$; if $post_loc = gf_prev_ptr$ then $print_ln(`.`)$ else print_ln(`,_after_special_info_at_byte_`, gf_prev_ptr : 1, `.`); $p \leftarrow signed_quad;$ if $p \neq gf_prev_ptr$ then $error(\texttt{backpointer_lin_byte_i}, cur_loc - 4: 1, \texttt{i_should_be_i}, gf_prev_ptr: 1, \texttt{i_not_i}, p: 1, \texttt{i});$ $design_size \leftarrow signed_quad; check_sum \leftarrow signed_quad;$ print(`design_size_i, design_size: 1, `_(`); print_scaled(design_size div 16); print_ln(`pt)`); $print_ln(\text{`check}_sum_l=l, check_sum : 1);$ $hppp \leftarrow signed_quad; vppp \leftarrow signed_quad;$ $print(\texttt{hppp}_{\square}\texttt{=}_{\square}\texttt{,}hppp:1,\texttt{`}_{\square}\texttt{(`)}; \ print_scaled(hppp); \ print_ln(\texttt{`)}\texttt{;} \ print(\texttt{`vppp}_{\square}\texttt{=}_{\square}\texttt{,}vppp:1,\texttt{`}_{\square}\texttt{(`)};$ $print_scaled(vppp); print_ln(`)`); pix_ratio \leftarrow (design_size/1048576) * (hppp/1048576);$ $min_m_stated \leftarrow signed_quad; max_m_stated \leftarrow signed_quad; min_n_stated \leftarrow signed_quad;$ $max_n_stated \leftarrow signed_quad;$ $print_ln(\min_{\square} m_{\square} = 1, min_m_stated : 1, max_{\square} m_{\square} = 1, max_m_stated : 1);$ if min_m_stated > min_m_overall then error(`min_m_should_be_<=`, min_m_overall: 1, `!`); if max_m_stated < max_m_overall then error(`max_m_should_be_>=`, max_m_overall: 1, `!`); $print_ln(\min_{\square}n_{\square}=_{\square}, min_n_stated: 1, \dots, max_{\square}n_{\square}=_{\square}, max_n_stated: 1);$ if min_n_stated > min_n_overall then error(`min_n_should_be_<=`, min_n_overall : 1, `!`); if $max_n_stated < max_n_overall$ then $error(\max_n_should_be_s=1, max_n_overall:1, 1, 1)$; $\langle Process the character locations in the postamble 65 \rangle;$ \langle Make sure that the end of the file is well-formed $64 \rangle$; end;

62. \langle Globals in the outer block 10 $\rangle +\equiv$ design_size, check_sum: integer; { TFM-oriented parameters } hppp, vppp: integer; { magnification-oriented parameters } post_loc: integer; { location of the post command } pix_ratio: real; { multiply by this to convert TFM width to scaled pixels }

63. (Set initial values 11) += $min_m_overall \leftarrow max_int; max_m_overall \leftarrow -max_int; min_n_overall \leftarrow max_int; max_n_overall \leftarrow -max_int;$ When we get to the present code, the *post_post* command has just been read.

64. \langle Make sure that the end of the file is well-formed 64 $\rangle \equiv$ if $k \neq post_post$ then $error(`should_be_postpost!`);$ for $k \leftarrow 0$ to 255 do if $char_ptr[k] > 0$ then $error(missing_|locator_|for_|character_|, k:1, '!');$ $q \leftarrow signed_quad;$ if $q \neq post_loc$ then $error(`postamble_pointer_should_be_`, post_loc: 1, `_not_`, q: 1, `!`);$ $m \leftarrow get_byte;$ if $m \neq gf_{id_byte}$ then $error(`identification_byte_should_be_`, gf_{id_byte} : 1, `, _not_`, m : 1, `!`);$ $k \leftarrow cur_loc; m \leftarrow 223;$ while $(m = 223) \land \neg eof(gf_file)$ do $m \leftarrow get_byte;$ if $\neg eof(gf_file)$ then $bad_gf(\text{`signature_in_byte_i'}, cur_loc - 1: 1, \text{`_should_be_223'})$ else if $cur_loc < k + 4$ then $error(int_enough_signature_bytes_at_end_of_file!)$ This code is used in section 61. \langle Process the character locations in the postamble $_{65}\rangle \equiv$ **65**. **repeat** $a \leftarrow cur_loc; k \leftarrow get_byte;$ if $(k = char_loc) \lor (k = char_loc + 1)$ then **begin** $c \leftarrow first_par(k);$ if $k = char_{-}loc$ then **begin** $u \leftarrow signed_quad$; $v \leftarrow signed_quad$; end else begin $u \leftarrow qet_byte * unity; v \leftarrow 0;$ end: $w \leftarrow signed_quad; p \leftarrow signed_quad; print(`Character_i, c:1, :_udx_i, u:1, :_u(');$ $print_scaled(u);$ if $v \neq 0$ then **begin** $print(), dy_{i}, v: 1, j(); print_scaled(v);$ end: $print(`), width_{u}, w: 1, `_{u}(`); w \leftarrow round(w * pix_ratio); print_scaled(w);$ $print_ln(), loc_l, p:1);$ if $char_ptr[c] = 0$ then $error(\text{`duplicate_llocator_lfor_this_character!`)}$ else if $p \neq char_ptr[c]$ then $error(character_location_should_be_', char_ptr[c]:1, '!');$ $char_ptr[c] \leftarrow 0; \ k \leftarrow no_op;$ end; until $k \neq no_{-}op$

This code is used in section 61.

66. The main program. Now we are ready to put it all together. This is where GFtype starts, and where it ends.

67. The main program needs a few global variables in order to do its work.

 $\langle \text{Globals in the outer block 10} \rangle +\equiv$ a: integer; { byte number of the current command } b, c, l, o, p, q, r: integer; { general purpose registers }

68. GFtype looks at the preamble in order to do error checking, and to display the introductory comment.

```
\langle \text{Process the preamble 68} \rangle \equiv
  open_gf_file; o \leftarrow get_byte; \{ fetch the first byte \}
  if o \neq pre then bad_gf(`First_byte_isn``t_start_of_preamble!`);
  o \leftarrow get_byte; \{ fetch the identification byte \} \}
  if o \neq gf_id_byte then bad_gf ('identification_byte_should_be_', gf_id_byte : 1, \text{`unot}, o: 1);
  o \leftarrow get_byte; \{ \text{fetch the length of the introductory comment} \}
  print(\cdots);
  while o > 0 do
     begin decr(o); print(xchr[get_byte]);
     end;
  print_ln(\cdots);
This code is used in section 66.
       \langle \text{Translate all the characters } 69 \rangle \equiv
69.
  repeat gf_prev_ptr \leftarrow cur_loc; (Pass no_op, xxx and yyy commands 70);
     if o \neq post then
        begin if o \neq boc then
          if o \neq boc1 then bad_gf(`byte_l`, cur_loc - 1: 1, `_lis_not_boc_l(`, o: 1, `)`);
        print_nl; print(cur_loc - 1: 1, ::\_beginning_of_char_j); \langle Pass a boc command 71 \rangle;
        if \neg do\_char then bad\_gf(\caruended\_unexpectedly\);
        max_n_observed \leftarrow n;
       if wants_pixels then \langle Print \text{ the image } 40 \rangle;
        \langle \text{Pass an } eoc \text{ command } 72 \rangle;
       end;
  until o = post;
```

This code is used in section 66.

```
70.
       \langle \text{Pass } no_op, xxx \text{ and } yyy \text{ commands } 70 \rangle \equiv
  repeat start_op;
     if o = yyy then
        begin \langle Translate a yyy command 55 \rangle;
        o \leftarrow no_{-}op;
        end
     else if (o \ge xxx1) \land (o \le xxx1 + 3) then
          begin (Translate an xxx command 53);
          o \leftarrow no_{-}op;
          end
        else if o = no_op then show_mnemonic(`no_op`);
  until o \neq no_{-}op;
This code is used in section 69.
71. (Pass a boc command 71) \equiv
  a \leftarrow cur\_loc - 1; incr(total\_chars);
  if o = boc then
     begin character_code \leftarrow signed_quad; p \leftarrow signed_quad; c \leftarrow character_code mod 256;
     if c < 0 then c \leftarrow c + 256;
     min\_m\_stated \leftarrow signed\_quad; max\_m\_stated \leftarrow signed\_quad; min\_n\_stated \leftarrow signed\_quad;
     max_n\_stated \leftarrow signed\_quad;
     \mathbf{end}
  else begin character_code \leftarrow get_byte; p \leftarrow -1; c \leftarrow character_code; q \leftarrow get_byte;
     max\_m\_stated \leftarrow get\_byte; min\_m\_stated \leftarrow max\_m\_stated - q; q \leftarrow get\_byte; max\_n\_stated \leftarrow get\_byte;
     min_n_stated \leftarrow max_n_stated - q;
     end;
  print(c:1);
  if character\_code \neq c then print(`\_with\_extension\_`, (character\_code - c) div 256 : 1);
  min_n_stated : 1, `<=n<=`, max_n_stated : 1);
  max_m_observed \leftarrow -1;
  if char_ptr[c] \neq p then
     error(\text{previous}_{\sqcup}\text{character}_{\sqcup}\text{pointer}_{\sqcup}\text{should}_{\sqcup}\text{be}_{\sqcup}, char_ptr[c]: 1, \text{,}_{\sqcup}\text{not}_{\sqcup}, p: 1, \text{,}')
  else if p > 0 then
        if wants_mnemonics then
           print_ln((previous_character_with_the_same_code_started_at_byte_(, p:1, ));
  char_ptr[c] \leftarrow qf_prev_ptr;
  if wants_mnemonics then print(`(initially_n=`, max_n_stated : 1, `)`);
  if wants_pixels then \langle \text{Clear the image } 38 \rangle;
  m \leftarrow 0; n \leftarrow 0; paint_switch \leftarrow white;
This code is used in section 69.
```

72. $\langle \text{Pass an } eoc \text{ command } 72 \rangle \equiv$

 $max_m_observed \leftarrow min_m_stated + max_m_observed + 1; n \leftarrow max_n_stated - max_n_observed; { now n is the minimum n observed }$

if $min_m_stated < min_m_overall$ then $min_m_overall \leftarrow min_m_stated$;

if $max_m_observed > max_m_overall$ then $max_m_overall \leftarrow max_m_observed$;

if $n < min_n$ -overall then min_n -overall $\leftarrow n$;

if $max_n_stated > max_n_overall$ then $max_n_overall \leftarrow max_n_stated$;

if $max_m_observed > max_m_stated$ then

 $print_ln(\text{`The}_previous_character_should_have_had_max_m_>=_`, max_m_observed : 1, `!`); if <math>n < min_n_stated$ then

 $print_ln($ The_previous_character_should_have_had_min_n<=_`, n: 1, `!`)

This code is used in section 69.

73. System-dependent changes. This section should be replaced, if necessary, by changes to the program that are necessary to make GFtype work at a particular installation. It is usually best to design your change file so that all changes to previous sections preserve the section numbering; then everybody's version will be consistent with the printed program. More extensive changes, which introduce new sections, can be inserted here; then only the index itself will get a new section number.

74. Index. Pointers to error messages appear here together with the section numbers where each identifier is used.

a: 24, 67abort: $\underline{7}$ aok: 49*print_scaled*: 45, 55, 61, 65b: 24, 67backpointer...should be p: 61backpointers: 18 $bad_char: 53, 54$ $bad_{-}gf: \underline{7}, 50, 64, 68, 69$ **while**: 0 **with**: 0 black: 14, 15, 35, <u>36</u>, 40, 57, 58, 59 *boc*: 13, 15, 16, 17, 18, 38, 42, 44, 48, 49, 52, 69, 71boc occurred before eoc: 52 boc1: 15, 16, 48, 52, 69break: 28proofing: 18ASCII_code: 8, 10, 27, 30buffer: 27, 29, 32, 33 byte n is not boc: 69 byte_file: 20, 21 $c: \underline{24}, \underline{30}, \underline{61}, \underline{67}$ **case**: 0 *decr*: $\underline{6}$, 43, 53, 68 char: 9char ended unexpectedly: 69 *char_loc*: 15, 16, 18, 48, 65 $char_loc0: 15$ *char_ptr*: 46, 47, 64, 65, 71character location should be...: 65 character_code: 46, 71check_sum: $61, \underline{62}$ cs: 17 $jump_out: 7$ cur_loc: 22, 23, 24, 50, 61, 64, 65, 69, 71 paint3: $\underline{15}$ $del_m: 15$ delta: 45design size: 17 $design_size: 61, 62$ *dialog*: 31, 66 $first_text_char: \underline{9}, 12$ **in**: 0 $do_{-}char:$ 44, 48, 49, 51, 69 xord: 10, 12, 29 $post: 13, 15, \underline{16}, 17, 19, 48, 52, 61, 62, 69$ $post_post: 15, 16, 17, 19, 48, 52, 64$ $ds: \underline{17}$ others: 2duplicate locator...: 65

dx: 15, 18xxx4: 15dy: 15, 18**type**: 0 $max_row: 5, 37, 38, 42$ eight_bits: <u>20</u>, 24, 48, 49 $eight_cases:$ 48 **div**: 0 else: 2 endcases: $\underline{2}$ eof: 24, 50, 64eoln: 29*error*: 50, 51, 52, 61, 64, 65, 71false: 36, 49, 53**then**: 0 chr: 9, 10, 12, 45final_end: 4, 7, 66First byte isn't...: 68 first_par: 48, 50, 51, 65**do**: 0 four_cases: 48, 51get: 29 $get_three_bytes: 24, 48$ $get_two_bytes: 24, 48$ $gf_{file}: \underline{3}, \underline{21}, 22, 23, 24, 50, 64$ $gf_id_byte: 15, 64, 68$ $gf_prev_ptr: 46, 61, 69, 71$ $GF_type: \underline{3}$ skip2: 15*i*: <u>3</u> identification byte should be n: 64, 68*image*: 37, 38, 43, 58 $image_array: 5, 37$ *incr*: $\underline{6}$, 24, 29, 38, 43, 53, 58, 59, 71 initialize: 3, 66input_ln: 27, <u>29</u>, 32, 33 integer: 3, 23, 24, 35, 39, 41, 45, 46, 48, 49,61, 62, 67 k: 29, 61l: 67lower_casify: $\underline{30}$, 32, 33*m*: $\underline{35}, \underline{61}$ n: 35wants_pixels: $\underline{25}$, 26, 33, 34, 57, 69, 71 $max_col: 5, 37, 38, 42$ $max_int: 63$ $max_m_observed: 40, 41, 42, 57, 71, 72$ $max_m_overall: 41, 61, 63, 72$ $max_m_stated: 38, 41, 61, 71, 72$ $max_n: 15, 17, 35, 38$

 $max_n_observed: 40, 41, 42, 69, 72$ $max_n_overall: 41, 61, 63, 72$ $max_n_stated: 38, \underline{41}, 43, 59, 60, 61, 71, 72$ $max_subcol:$ 38, <u>39</u>, 40, 42, 43, 58 $max_subrow: 38, 39, 42, 43, 58$ **set**: 0 $new_row_0: 15, 16, 48, 51$ check sum: 17 Chinese characters: 18 $min_m: 15, 17, 35, 38$ $min_m_overall: 41, 61, 63, 72$ $min_m_stated: 38, 41, 43, 61, 71, 72$ $min_n_overall: 41, 61, 63, 72$ $min_n_stated: 38, 41, 61, 71, 72$ missing locator...: 64 dm: 15Mnemonic output?: 32 *eoc*: 13, 15, $\underline{16}$, 17, 48, 52 **not**: 0 *hppp*: 17, 61, 62 $print_ln: \underline{3}, \underline{34}, 40, 42, 43, 49, 61, 65, 68, 71, 72$ Fuchs, David Raymond: 1, 19 $last_text_char: \underline{9}, 12$ *negate*: $\underline{6}$, $\underline{45}$ $new_row_1: 15$ $new_row_164: 15$ $text_char: 9, 10$ nl_{error} : 50, 53 non-ASCII character...: 53 not enough signature bytes...: 64 o: 49, 67*line_length*: 5, 53boolean: 25, 36, 49, 54 $open_gf_file: \underline{22}, 68$ Options selected: 34 ord: 10, 45oriental characters: 18 othercases: 2 xxx1: 15, <u>16</u>, 48, 51, 70 $p: \underline{49}, \underline{61}, \underline{67}$ x: 0Bad GF file: 7 paint: 56 $paint_0: 15, 16, 48$ paint1: 15, 16, 48, 51, 56paint2: 15*reset*: 22, 29 $get_byte: 24, 48, 50, 53, 64, 65, 68, 71$ $min_n: 15, 17, 38$ *pix_ratio*: 61, 62, 65*pixel*: 35, 36, 37Pixel output?: 33

to: 0 $no_op:$ 15, <u>16</u>, 18, 48, 52, 65, 70 *post_loc:* 61, 62, 64postamble command within...: 52 postamble pointer should be...: 64 Postamble starts at byte n: 61 *pre*: 13, 15, $\underline{16}$, 48, 52, 68 preamble command within...: 52 previous character...: 71, 72 *print*: $\underline{3}$, 7, 34, 43, 45, 50, 53, 55, 56, 57, 59, 60, 61, 65, 66, 68, 69, 71 $print_nl: \underline{3}, 43, 50, 52, 53, 66, 69$ function: 0 output: 3 q: 49, 61, 67banner: $\underline{1}, 3, 31$ read: 24 $read_ln: 29$ $read_postamble: \underline{61}, \underline{66}$ real: 62*rewrite*: 31round: 6524d: r: 67*s*: 45 y: 0scaled: 15, 17, 18 record: 0should be postpost: 64 show_label: 50 show_mnemonic: <u>50</u>, 52, 53, 55, 59, 60, 70 signature...should be...: 64 signed_quad: 24, 48, 61, 64, 65, 71sixteen_cases: 48 $sixty_four_cases: 48, 51$ *skip0*: 15, <u>16</u>, 48, 51 skip1: 15, 16, 48, 60 *unity*: 45, 65start_op: 50, 56, 70 string of negative length: 53 system dependencies: 2, 7, 9, 19, 20, 24, 25, 27, 28, 29, 31, 36, 37, 38, 40, 73 Japanese characters: 18 term_in: 27, 29 $term_out: 27, 28, 31, 32, 33$ $terminal_line_length: 5, 27, 29$ text_file: 9, 27The character is too large...: 42the file ended prematurely: 50The file had n characters...: 66 thirty_seven_cases: 48, 51 $thirty_two_cases: 48$

This pixel's lower...: 43This pixel's upper: 43 *skip3*: 15 **end**: 2 Knuth, Donald Ervin: 1 $total_chars: \ \underline{46}, \ 47, \ 66, \ 71$ **or**: 0 true: 25, 26, 36, 49, 53 *xxx2*: 15 $u: \underline{61}$ $max_m: 15, 17, 38$ undefined command: 51 $undefined_commands: 16, 48$ $update_terminal: 28, 29$ v: 61paint_switch: <u>14</u>, 15, <u>35</u>, 57, 58, 59, 60, 71 $del_n: \underline{15}$ vppp: 17, 61, 62*w*: 61wants_mnemonics: 25, 26, 32, 34, 50, 53, 55, 56, 57, 59, 60, 71 **if**: 0 white: $15, 35, \underline{36}, 38, 40, 43, 57, 58, 60, 71$ write: 3, 32, 33 write_ln: 3, 31, 32, 33*xchr*: 10, 11, 12, 53, 68**file**: 0 yyy: 15, <u>16</u>, 18, 48, 51, 70

- $\langle \text{Cases for commands } no_op, pre, post, post_post, boc, and eoc 52 \rangle$ Used in section 51.
- (Clear the image 38) Used in section 71.
- $\langle \text{Compare the subarray boundaries with the observed boundaries 42} \rangle$ Used in section 40.
- $\langle \text{Constants in the outer block 5} \rangle$ Used in section 3.
- (Determine whether the user *wants_mnemonics* 32) Used in section 31.
- $\langle \text{Determine whether the user wants_pixels 33} \rangle$ Used in section 31.
- (Globals in the outer block 10, 21, 23, 25, 27, 35, 37, 39, 41, 46, 54, 62, 67) Used in section 3.
- $\langle \text{Labels in the outer block 4} \rangle$ Used in section 3.
- \langle Make sure that the end of the file is well-formed 64 \rangle Used in section 61.
- (Paint pixels m p through m 1 in row n of the subarray 58) Used in section 57.
- (Paint the next p pixels 57) Used in section 56.
- $\langle \text{Pass a boc command } 71 \rangle$ Used in section 69.
- (Pass an *eoc* command 72) Used in section 69.
- $\langle Pass no_op, xxx and yyy commands 70 \rangle$ Used in section 69.
- $\langle Print all the selected options 34 \rangle$ Used in section 31.
- (Print asterisk patterns for rows 0 to $max_subrow 43$) Used in section 40.
- $\langle Print the image 40 \rangle$ Used in section 69.
- $\langle Process the character locations in the postamble 65 \rangle$ Used in section 61.
- $\langle Process the preamble 68 \rangle$ Used in section 66.
- $\langle \text{Set initial values } 11, 12, 26, 47, 63 \rangle$ Used in section 3.
- \langle Start translation of command o and **goto** the appropriate label to finish the job 51 \rangle Used in section 50.
- \langle Translate a sequence of *paint* commands, until reaching a non-*paint* 56 \rangle Used in section 51.
- $\langle \text{Translate a } new_row \text{ command } 59 \rangle$ Used in section 51.
- $\langle \text{Translate a } skip \text{ command } 60 \rangle$ Used in section 51.
- $\langle \text{Translate a } yyy \text{ command } 55 \rangle$ Used in sections 51 and 70.
- $\langle \text{Translate all the characters } 69 \rangle$ Used in section 66.
- $\langle \text{Translate an } xxx \text{ command } 53 \rangle$ Used in sections 51 and 70.
- (Translate the next command in the GF file; goto 9999 if it was eoc; goto 9998 if premature termination is needed 50) Used in section 49.
- \langle Types in the outer block 8, 9, 20, 36 \rangle Used in section 3.